A linear-time algorithm to sample the dual-birth model

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Abstract—The ability to sample models of tree evolution is essential in the analysis and interpretation of phylogenetic trees. The dual-birth model is an extension of the traditional birth-only model and allows for sampling trees of varying degrees of balance. However, for a tree with \( n \) leaves, the tree sampling algorithm proposed in the original paper is \( O(n \log n) \). I propose an algorithm to sample trees under the dual-birth model in \( O(n) \), and I provide a fast C++ implementation of the proposed algorithm.

I. INTRODUCTION

The dual-birth model is a birth-only model in which each branch is a Poisson process with one of two states: active, in which splits occur with rate \( \lambda_a \), and inactive, in which splits occur with rate \( \lambda_b \). It is typically parameterized by \( \lambda = \lambda_a + \lambda_b \) and \( r = \lambda_a/\lambda_b \) [1]. When a branch splits, one child branch is active and the other is inactive. Simulating trees under the dual-birth model can provide null splits, one child branch is active and the other is inactive. Because each branch is a Poisson process, the time until the next split is an exponential random variable with the rate of splitting event is the minimum of \( n_a \) exponential random variables, each with rate \( \lambda_a \), and \( n_i \) exponential random variables, each with rate \( \lambda_i \).

The minimum of \( k \) exponential random variables \( X_1, \ldots, X_k \) with rates \( \lambda_1, \ldots, \lambda_k \) is an exponential random variable with rate \( \sum_{i=1}^{k} \lambda_i \). Therefore, the time until the next splitting event is an exponential random variable with rate \( n_a \lambda_b + n_i \lambda_a \).

Given \( k \) exponential random variables \( X_1, \ldots, X_k \) with rates \( \lambda_1, \ldots, \lambda_k \), the probability that \( X_i \) is the minimum of all \( k \) random variables is \( \lambda_i / \sum_{i=1}^{k} \lambda_i \). Because only one terminal branch can be the next to split, the events that pendant branch \( j \) (for all \( j \)) is the next to split are mutually exclusive. Therefore, the probability that the next branch to split is an active branch is \( p = n_a \lambda_b / (n_a \lambda_b + n_i \lambda_a) \). Once the state of the next branch to split is determined (by sampling a Bernoulli random variable with probability \( p \)), because all leaves in that state have the same rate (and are thus equally likely to be the next to split), the next splitting leaf can be chosen uniformly from the set of leaves in the chosen state.

II. METHODS

A. Linear-time algorithm description

Let \( \exp(\lambda) \) denote an exponential random variable with rate \( \lambda \). Let \( \text{bern}(p) \) denote an exponential random variable with probability \( p \). Let \( \| x \| \) denote the size of set \( x \). The algorithm starts with two sets of leaves, active (with just the root) and inactive (initially empty). The algorithm iterates while the termination condition has not yet been met. The time until the next split is sampled from \( \exp(\lambda_b \| \text{active} \| + \lambda_a \| \text{inactive} \|) \). The next splitting leaf is chosen by first determining its state by sampling \( \text{bern}(\lambda_b \| \text{active} \| / (\lambda_b \| \text{active} \| + \lambda_a \| \text{inactive} \|)) \) (success = active, failure = inactive) and then uniformly choosing a random leaf from the corresponding set. The chosen leaf’s time is set to the sampled split time, the leaf is removed from its set, two children are created, and one child is placed in active and the other in inactive. When the termination condition has been met, the times of all remaining leaves in active and inactive are set to the end time.

B. Algorithm correctness

The algorithm begins with a single active leaf at time 0. Because each branch is a Poisson process, the time until the branch splits is an exponential random variable with the branch’s rate. Therefore, at a given time \( t \) with \( n_a \) active leaves and \( n_i \) inactive leaves, the time until the next splitting event is the minimum of \( n_a \) exponential random variables, each with rate \( \lambda_a \), and \( n_i \) exponential random variables, each with rate \( \lambda_i \).

My implementation of the proposed algorithm closely follows the theoretical expectations of the dual-birth model (Fig. 1) and is orders of magnitude faster than the original implementation of the initially-proposed algorithm (Fig. 2).

III. RESULTS

My implementation of the proposed algorithm scales linearly with the number of leaves, and its C++ implementation is fast and does not have any non-STL dependencies, making it easy to build on all environments. My hope is that the fast implementation I provide can be wrapped into existing tools for the convenience of users as well as developers.

IV. DISCUSSION

The proposed algorithm scales linearly with the number of leaves, and its C++ implementation is fast and does not have any non-STL dependencies, making it easy to build on all environments. My hope is that the fast implementation I provide can be wrapped into existing tools for the convenience of users as well as developers.
Theoretical expectations of (a) cherry fraction (dashed red line) and active leaf fraction (dashed purple line), and (b) branch length (dashed blue line) and terminal branch length (dashed green line) versus simulated distributions (in box plots) using 100 replicates with $n = 4096$, $\lambda = 48$, and varying values of $r$ (x-axis) from $1/1024$ to $1$.

Fig. 1. Theoretical expectations of (a) cherry fraction (dashed red line) and active leaf fraction (dashed purple line), and (b) branch length (dashed blue line) and terminal branch length (dashed green line) versus simulated distributions (in box plots) using 100 replicates with $n = 4096$, $\lambda = 48$, and varying values of $r$ (x-axis) from $1/1024$ to $1$.

Fig. 2. Log-scaled runtimes of the original Python-implemented algorithm and the new C++-implemented algorithm for $1000 \leq n \leq 10000$ for $r = 10^{-4}$ and $r = 1$. Each measured time is the average of the time taken to simulate 20 trees in a single execution, and each point plot is the distribution of 20 replicate executions.

**AVAILABILITY**

My C++ implementation of the proposed algorithm can be found in the following GitHub repository: [https://github.com/niemasd/Dual-Birth-Simulator](https://github.com/niemasd/Dual-Birth-Simulator)

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**REFERENCES**